



ANSWER BOOK

Dynamic programming

Decision Mathematics 2 · Year FM

6 questions · 21 marks

NAVY EDITION

SECTION A · Warm-up

- A1** Any part of an optimal path is itself optimal between its own two ends. It allows the problem to be solved **backwards**: once the best value from every vertex onwards is known, the best value at the vertex before it is one comparison away, and no route already rejected needs reconsidering. [2]
- A2** Minimum total: $\min(6 + 11, 10 + 5) = \min(17, 15) = 15$, through the second option.
Minimax: $\min(\max(6, 11), \max(10, 5)) = \min(11, 10) = 10$, also through the second, but chosen because its worst single stage is smaller rather than because its total is. [2]

SECTION B · Standard

- B1** Stage 3: from C the value is **6**, from D it is **7**.
Stage 2: A gives $\min(4 + 6, 9 + 7) = \min(10, 16) = 10$ through C; B gives $\min(8 + 6, 2 + 7) = \min(14, 9) = 9$ through D.
Stage 1: S gives $\min(5 + 10, 3 + 9) = \min(15, 12) = 12$ through B.
The route is **S, B, D, T** of length 12. Checking all four routes gives 12, 15, 17 and 21, confirming it. [6]
- B2** The two happen to agree here, which is exactly why the example is a poor test of the method. The choice was made by comparing $5 + 10$ with $3 + 9$, both of which use values found from the far end. Had the value at A been 5 rather than 10, S to A would have won at 10 despite the dearer first step. A greedy first move carries no guarantee about what follows it. [3]
- B3** The **stage** says how far through the process you are: here there are three stages, one per arc travelled. The **state** says where you are within that stage: at stage 2 the states are A and B, at stage 3 they are C and D. Every row of the table is one (stage, state) pair, and naming both is part of the answer. [3]

SECTION C · Stretch

- C1** This is a **maximin** problem, so the addition is replaced by taking the minimum along a route.
 Stage 3: C is worth 6 and D is worth 7, as before.
 Stage 2: A gives $\max(\min(4, 6), \min(9, 7)) = \max(4, 7) = 7$ through D; B gives $\max(\min(8, 6), \min(2, 7)) = \max(6, 2) = 6$ through C.
 Stage 1: S gives $\max(\min(5, 7), \min(3, 6)) = \max(5, 3) = 5$ through A.
 The route is **S, A, D, T** with smallest stage **5**. It is a different route from the shortest one, which shows that the two objectives genuinely pull apart: S, A, D, T is the longest route of the four by total length, at 21.
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