



ANSWER BOOK

The Simplex algorithm

Decision Mathematics 1 · Year FM

6 questions · 20 marks

NAVY EDITION

SECTION A · Warm-up

- A1** Stop when there are no negative entries left in the objective row. A negative entry [2]
means increasing that variable would still improve the objective, so while one remains
the solution can be bettered. Once none remain, no variable can be increased without
reducing the objective, so the current corner is optimal.
- A2** Add slack variables: $x + y + s_1 = 10$ and $2x + 3y + s_2 = 24$, with the objective as $P - 4x - 5y = 0$. The tableau rows are $P: 1, -4, -5, 0, 0 \mid 0$; $s_1: 0, 1, 1, 1, 0 \mid 10$; $s_2: 0, 2, 3, 0, 1 \mid 24$.

SECTION B · Standard

- B1** The most negative entry in the objective row is -5 , so the pivot column is y [4]
Ratio test: $10 \div 1 = 10$ and $24 \div 3 = 8$, so the pivot row is the s_2 row. Divide that row by 3, then subtract it from the s_1 row and add
five times it to the objective row. That gives $P = 40$ with $y = 8$, and the
objective row still has a negative entry in x .
- B2** No entry in the objective row is negative, so no variable can be increased to improve P [4]
the tableau is optimal. The basic variables are $x = 6$ and $y = 4$, the non-basic slacks are $s_1 = s_2 = 0$,
and $P = 44$. Both slacks being zero means both constraints are tight,
matching the graphical answer.
- B3** Ratios: $12 \div 2 = 6$ and $20 \div 4 = 5$. The -1 is skipped [3]
entirely: increasing the entering variable would make that row's value grow rather than
shrink, so it places no limit. The smallest ratio is 5, so the third row is
the pivot row.

SECTION C · Stretch

- C1** With only $\leq$ constraints the origin is feasible and the slack variables form a valid starting [4]
solution. A $\geq$ constraint excludes the origin, so setting all the real variables to zero gives
a negative surplus and no starting corner. The two-stage method
first minimises the sum of the artificial variables, driving them to zero and so reaching a
feasible corner, then runs ordinary Simplex from there. The big-M method
instead subtracts a large multiple M of each artificial variable from the
objective, making them so expensive that the algorithm eliminates them on its own.